

# Indexing and Searching

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## References:

1. Modern Information Retrieval, chapter 9
2. Information Retrieval: Data Structures & Algorithms, chapter 5
3. [G.H. Gonnet, R.A. Baeza-Yates, T. Snider, Lexicographical Indices for Text: Inverted files vs. PAT trees](#)

# Introduction

- The main purpose of an IR system is
  - To help users find information of their interest, achieving high **effectiveness** (maximizing the ratio of user satisfaction versus user effort)
- Here we look at the other side of the coin
  - Viz. the secondary issue, **efficiency**
    - To process user queries with minimal requirements of computational resources, network bandwidth, etc.
    - As we move to larger-scale applications, efficiency becomes more and more important

# Introduction (cont.)

- **Sequential or online searching**

- Find the occurrences of a pattern in a text when the **text is not preprocessed**
  - Only appropriate when:
    - The text is small
    - Or the text collection is **very volatile**
    - Or the index space overhead cannot be afforded

- **Indexed search**

- Build data structures over the text (i.e., indices) to speed up the search
- Appropriate for the larger or semi-static text collection
- The system updated at reasonably regular intervals

# Introduction (cont.)

- The efficiency of an indexed IR system can be measured by:
  - Indexing time:
    - The time needed to build the index
  - Indexing space:
    - Space used during the generation of the index
  - Index storage:
    - Space required to store the index
  - Query latency:
    - Time interval between the arrival of the query in the IR system and the generation of the answer
  - Query throughput:
    - Average number of query processed per second.

# Introduction (cont.)

- Three **data structures** for indexing are considered

- **Inverted files**

- The best choice for most applications

- **Signature files**

- Popular in the 1980s

- **Suffix arrays**

- Faster but harder to build and maintain

Issues:

Search cost,  
Space overhead,  
Building/updating time

# Inverted Files

- A word-oriented mechanism for indexing a text collection in order to speed up the searching task
  - Two elements:
    - A vector containing all the distinct words (called **vocabulary**) in the text collection
      - The space required for the vocabulary is rather small:
        - $\sim O(n^\beta)$ ,  $n$ : the text size,  $0 < \beta < 1$  (Heaps' law)
    - For each vocabulary word, a list of all docs (**identified by doc number in ascending order**) in which that word occurs
      - Space overhead: 30~40% of the text size (for text position addressing)
- Distinction between inverted file or list
  - **Inverted file**: occurrence points to documents or file names (identities)
  - **Inverted list**: occurrence points to word positions

# Inverted Files (cont.)

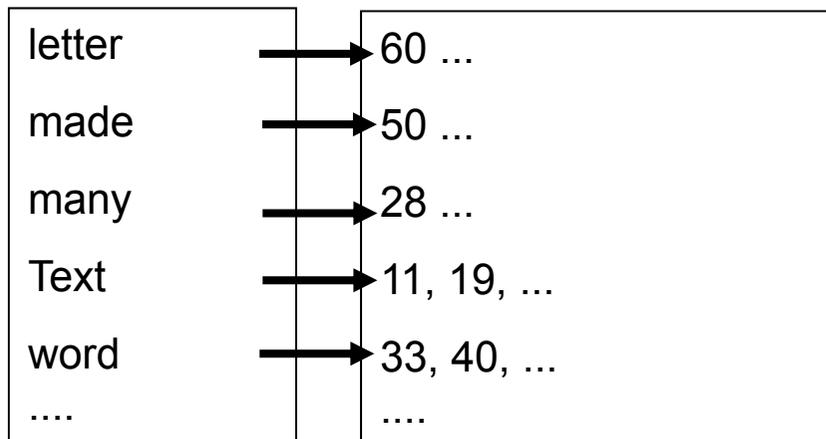
- **Example**

1    6   9 11   17 19   24 28   33    40    46 50   55   60  
This is a text. A text has many words. Words are made from letters.

**Text**

**Vocabulary**

**Occurrences**



An inverted list

Each element in a list points to a text position

An inverted file

Each element in a list points to a doc number

*difference:  
indexing granularity*

# Inverted Files: Addressing Granularity

- Text (word/character) positions (full inverted indices)
- Documents
  - All the occurrences of a word inside a document are collapsed to one reference
- (Logical) blocks
  - The blocks can be of fixed or different size
  - All the occurrences of a word inside a single block are collapsed to one reference
  - Space overhead: ~5% of the text size for a large collection

# Inverted Files: Some Statistics

- Size of an inverted file as approximate percentages of the size of the text collection

Index		Small Collection (1 Mb)		Medium Collection (200 Mb)		Large Collection (2 Gb)	
4 bytes/pointer	Addressing Words	45%	73%	36%	64%	35%	63%
1,2,3 bytes/pointer	Addressing Documents	19%	26%	18%	32%	26%	47%
2 bytes/pointer	Addressing 64K blocks	27%	41%	18%	32%	5%	9%
1 byte/pointer	Addressing 256 blocks	18%	25%	1.7%	2.4%	0.5%	0.7%

Stopwords are removed

Stopwords are indexed

# Inverted Files (cont.)

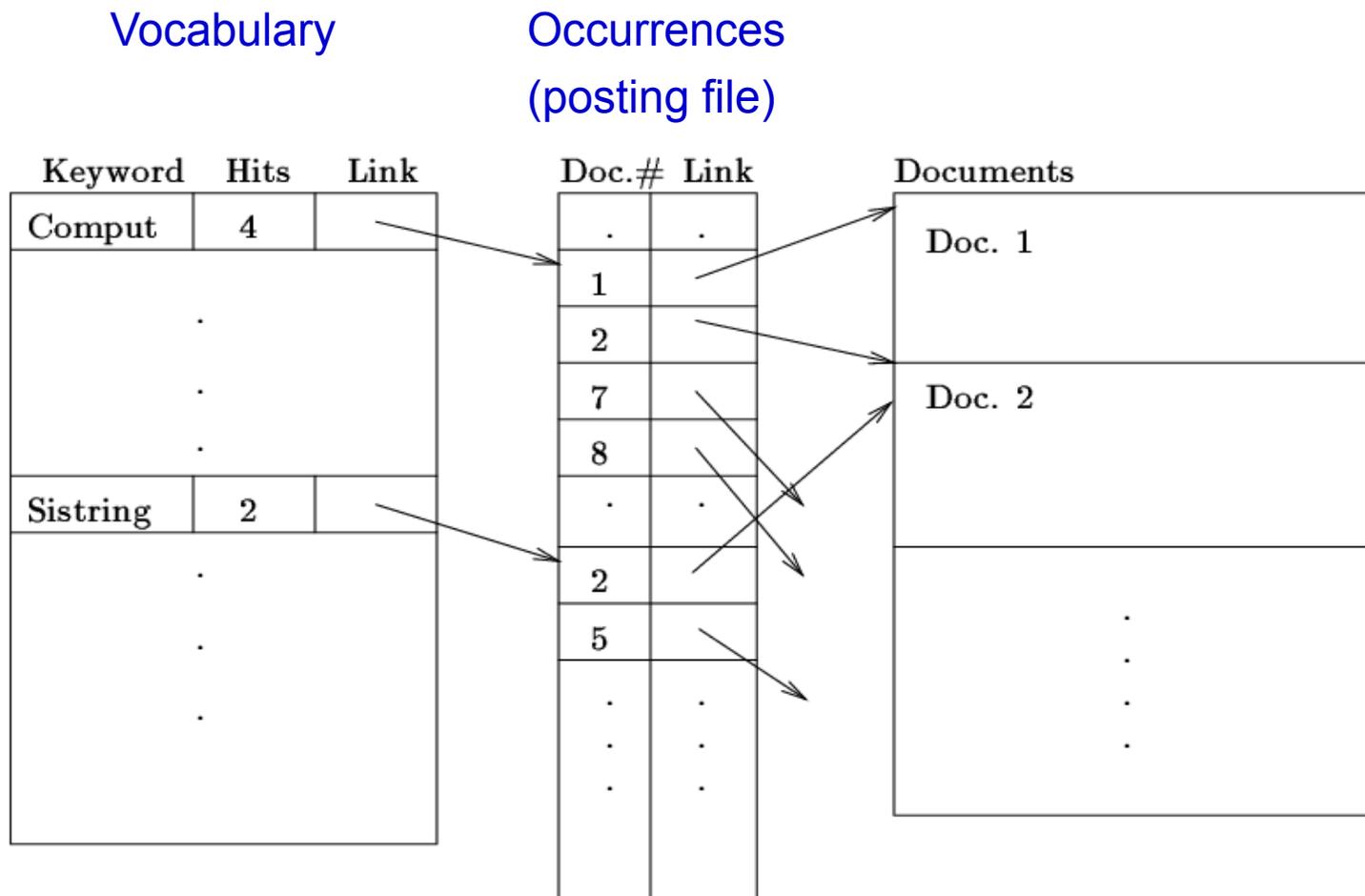
- **Document addressing**

- Assume that the vocabulary (control dictionary) can be kept in main memory. Assign a sequential word number to each word
- Scan the text database and output to a temporary file containing the record number and its word number
- Sort the temporary file by word number and use record number as a minor sorting field
- Compact the sorted file by removing the word number. During this compaction, build the inverted list from the end points of each word. This compacted file (**posting file**) becomes the main index

.....  
 $d_5 w_3$   
 $d_5 w_{100}$   
 $d_5 w_{1050}$   
.....  
 $d_9 w_{12}$   
.....

# Inverted Files (cont.)

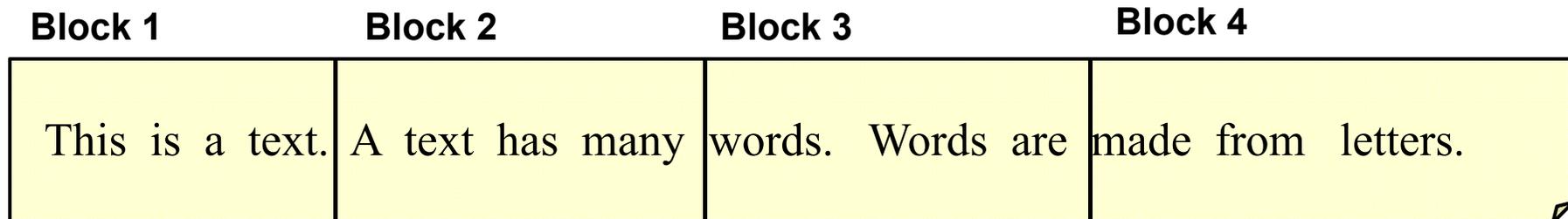
- **Document addressing (count.)**



# Inverted Files: Block Addressing

- Features
  - Text is divided into blocks
  - The occurrences in the invert file point to blocks where the words appear
  - Reduce the space requirements for recording occurrences
- Disadvantages
  - The occurrences of a word inside a single block are collapsed to one reference
  - Online search over qualifying blocks is needed if we want to know the exact occurrence positions
    - Because many **retrieval units** are packed into a single block

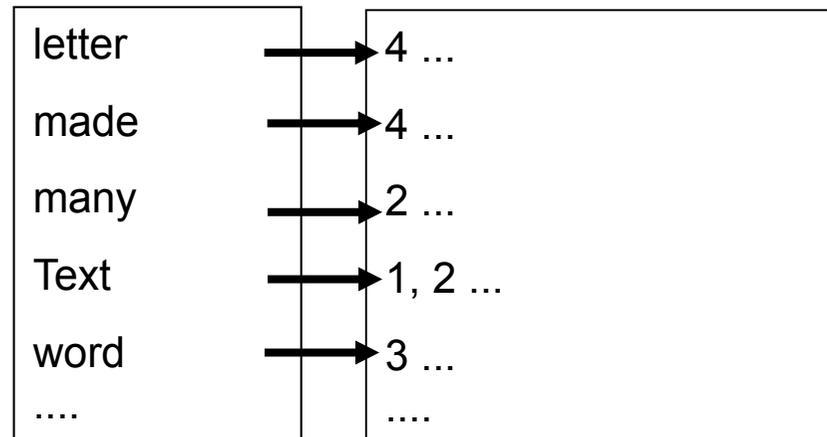
# Inverted Files: Block Addressing (cont.)



**Text**

**Vocabulary**

**Occurrences**



**Inverted Index**

# Inverted Files: Searching

- Three general steps
  - **Vocabulary search**
    - Words and patterns in the query are isolated and searched in the vocabulary
    - Phrase and proximity queries are split into single words  
"white house"                      "network of computer"    "computer network"
  - **Retrieval of occurrences**
    - The lists of the occurrences of all words found are retrieved
  - **Manipulation of occurrences**                      *intersection, distance, etc.*
    - For phrase, proximity or Boolean operations
    - Directly search the text if block addressing is adopted

# Inverted Files: Searching (cont.)

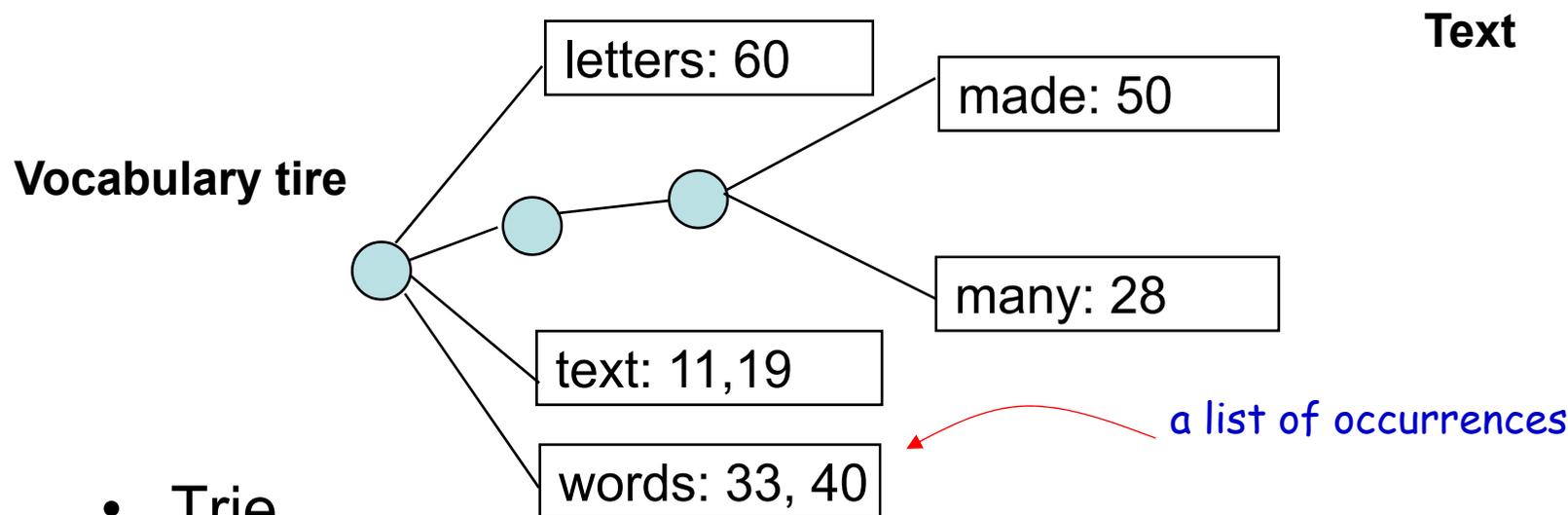
- Most time-demanding operation on inverted files is **the merging or intersection** of the lists of occurrences
  - E.g., **for the context queries**
    - Each element (word) searched separately and a list (occurrences for word positions, doc IDs, ..) generated for each
- The lists of all elements traversed in synchronization to find places where all elements appear in sequence (for a phrase) or appear close enough (for proximity)

*An expansive solution*

# Inverted Files: Construction

- The **trie** data structure to store the **vocabulary**

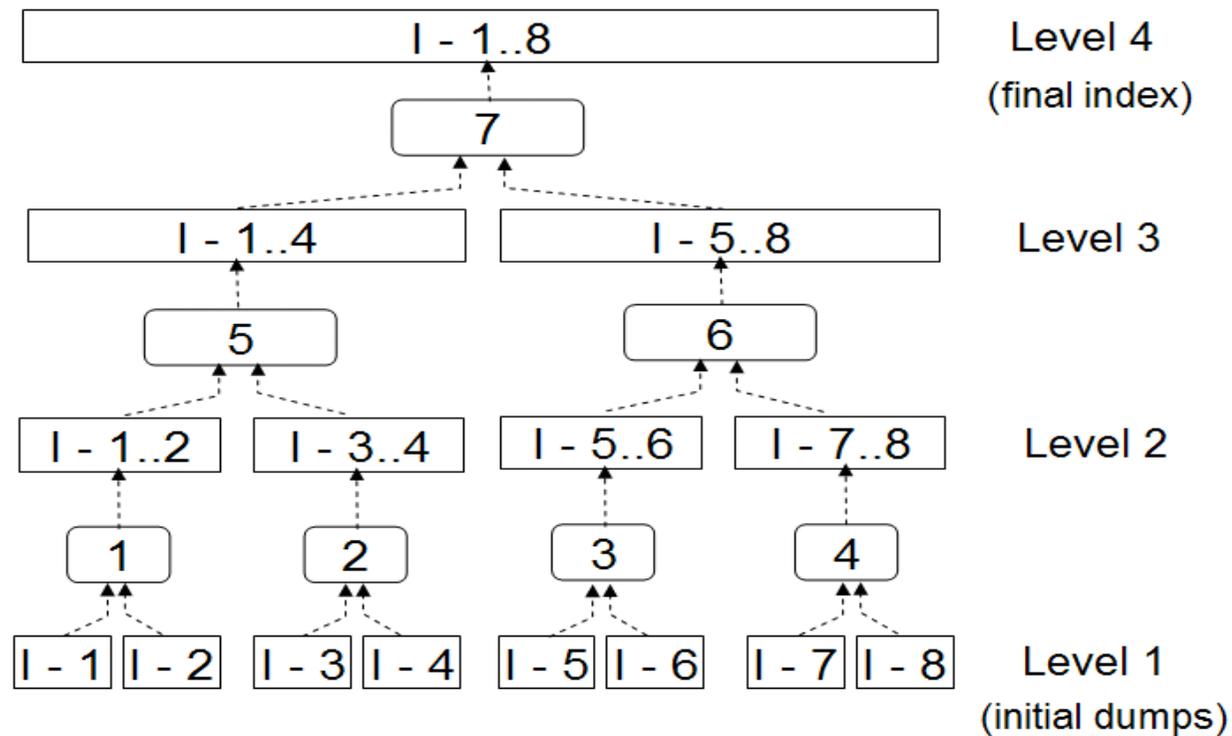
1 6 9 11 17 19 24 28 33 40 46 50 55 60  
This is a text. A text has many words. Words are made from letters.



- **Trie**
  - A digital search tree
  - A **multiway tree** that stores set of strings and able to retrieve any string in time proportional to its length
  - A special character is added to the end of string to ensure that no string is a prefix of another (words appear only at leaf nodes)

# Inverted Files: Construction (cont.)

- Merging of the partial indices
  - Merge the sorted vocabularies
  - Merge both lists of occurrences if a word appears in both indices



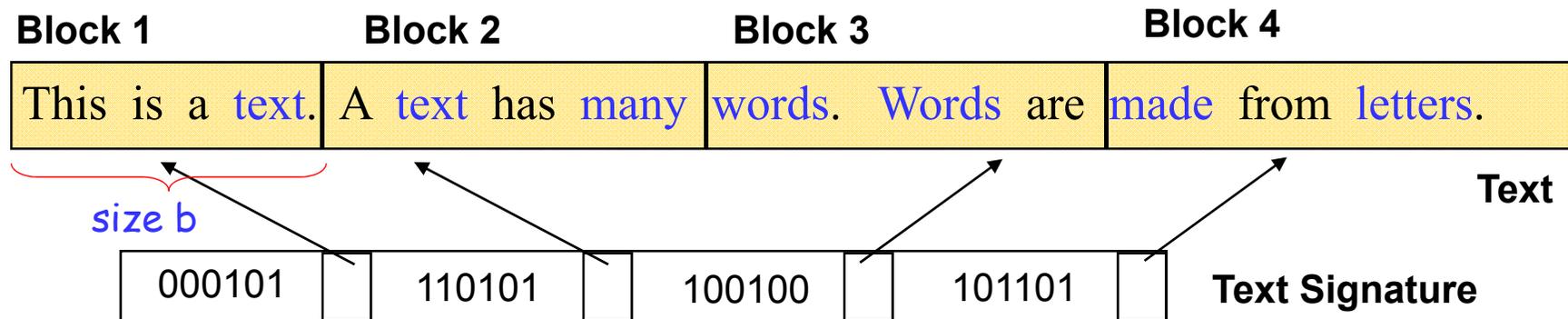
# Inverted Files: Performance

- For a full index built on 250 Mb of text
  - Single word: 0.08 sec
  - Phrase (2~5 words): 0.25 to 0.35 sec

# Signature Files

- **Basic Ideas**
  - **Word-oriented index structures based on hashing**
    - A hash function (signature) maps words to bit masks of  $B$  bits
  - Divide the text into **blocks of  $b$  words** each
    - **A bit mask of  $B$  bits** is assigned to each block by **bitwise ORing the signatures of all the words in the text block**
  - A word is presented in a text block if all bits set in its signature are also set in the bit mask of the text block

# Signature Files (cont.)



## Signature functions

$h(\text{text})$	=	000101
$h(\text{many})$	=	110000
$h(\text{words})$	=	100100
$h(\text{made})$	=	001100
$h(\text{letters})$	=	100001

size B

## Stop word list

this
is
a
has
are
from
.....

- The text signature contains
  - Sequences of bit masks
  - Pointers to blocks

# Signature Files (cont.)

- **False Drops** or False Alarms
  - All the corresponding bits are set in the bit mask of a text block, but the query word is not there
  - E.g., a false drop for the index “letters” in block 2
- **Goals** of the design of signature files
  - Ensure the probability of a false drop is low enough
  - Keep the signature file as short as possible

tradeoff

# Signature Files: Searching

- **Single word queries**

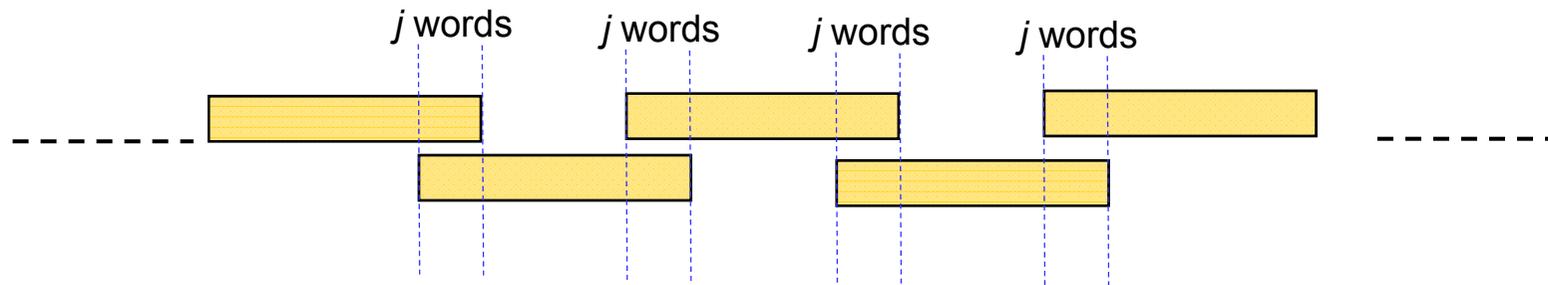
- Hash each word to a bit mask  $W$
- Compare the bit mask  $B_i$  of all text block (linear search) if they contain the word ( $W \& B_i == W$  ?)
  - **Overhead:** online traverse candidate blocks to verify if the word is actually there

- **Phrase or Proximity queries**

- The bitwise OR of all the query (word) masks is searched
- The candidate blocks should have the same bits presented “1” as that in the composite query mask
- **Block boundaries should be taken care of**
  - For phrases/proximities across two blocks

# Signature Files: Searching (cont.)

- **Overlapping blocks**



- **Other types of patterns** (e.g., prefix/suffix strings,...) are not supported for searching in this scheme

- **Construction**

- Text is cut in blocks, and for each block an entry of the signature file is generated
  - Bitwise OR of the signatures of all the words in it
- Adding text and deleting text are easy

# Signature Files: Searching (cont.)

- **Pros**

- Pose a low overhead (10-20% text size) for the construction of text signature
- Efficient to search phrases and reasonable proximity queries (**the only scheme improving the phrase search**)

- **Cons**

- Only applicable to index words
- Only suitable for not very large texts
  - Sequential search (check) in the text blocks to avoid false drops
  - Inverted files outperform signature files for most applications

# Signature Files: Performance

- For a signature file built on 250 Mb of text
  - Single word (or phrase?): 12 sec

# Suffix Trees

- **Premise**

- Inverted files or signature files treat the text as a sequence of words
  - For collections that the concept of word does not exist, they would be not feasible (like genetic databases)

- **Basic Ideas**

- Each position (character or bit) in the text considered as a text suffix
  - A string going from that text position to the end of the text (arbitrarily far to the right)
- Each suffix (or called **semi-infinite string**, *sistring*) uniquely identified by its position
  - Two suffixes at different positions are lexicographical different

A special null character is added to the strings' ends

# Suffix Trees (cont.)

- **Basic Ideas (cont.)**

- **Index points:** not all text positions indexed
  - Word beginnings
  - Or, beginnings of retrievable text positions
- Queries are based on prefixes of sistrings, i.e., on any substring of the text

1	6	9	11	17	19	24	28	33	40	46	50	55	60
This is a <b>text</b> . A <b>text</b> has <b>many words</b> . <b>Words</b> are <b>made</b> from <b>letters</b> .													

- sistring 11: text. A text has many words. Words are made from letters.
- sistring 19: text has many words. Words are made from letters.
- sistring 28: many words. Words are made from letters.
- sistring 33: words. Words are made from letters.
- sistring 40: Words are made from letters.
- sistring 50: made from letters.
- sistring 60: letters.

# Suffix Trees (cont.)

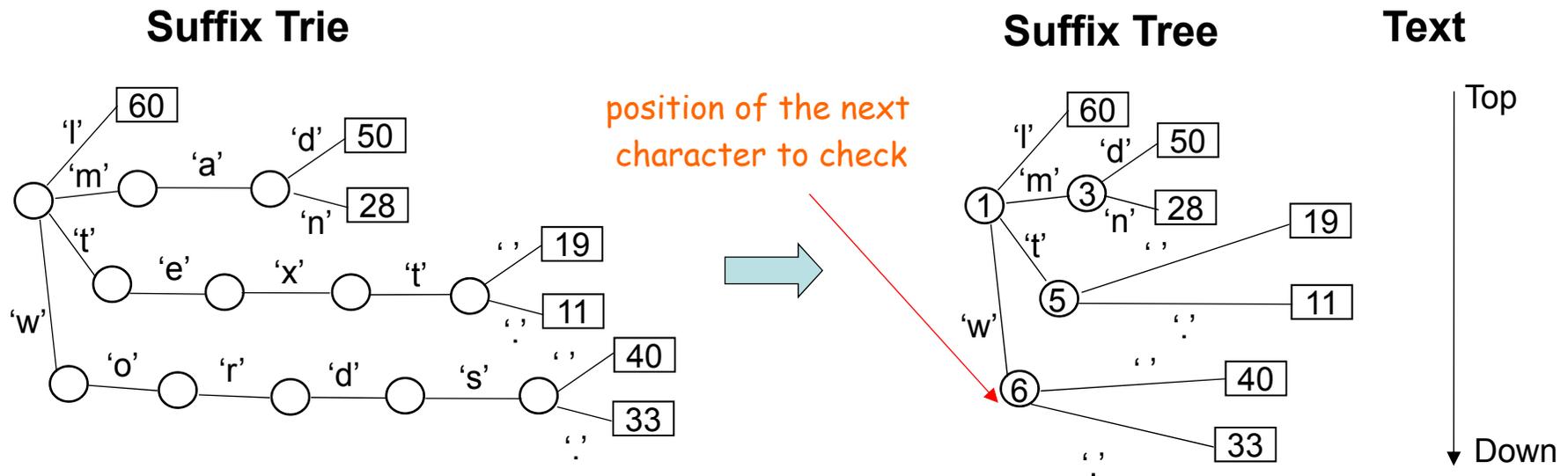
- **Structure**

- The suffix tree is a trie structure built over all the suffixes of the text
  - Points to text are stored at the leaf nodes
- The suffix tree is implemented as a **Patricia tree** (or **PAT tree**), i.e., **a compact suffix tree**
  - Unary paths (where each node has just one child) are compressed
  - An indication of next character (or bit) position to consider/check are stored at the internal nodes
    - Each node takes 12 to 24 bytes
    - A space overhead of 120%~240% over the text size

# Suffix Trees (cont.)

- PAT tree over a sequence of characters

1 6 9 11 17 19 24 28 33 40 46 50 55 60  
 This is a text. A text has many words. Words are made from letters.



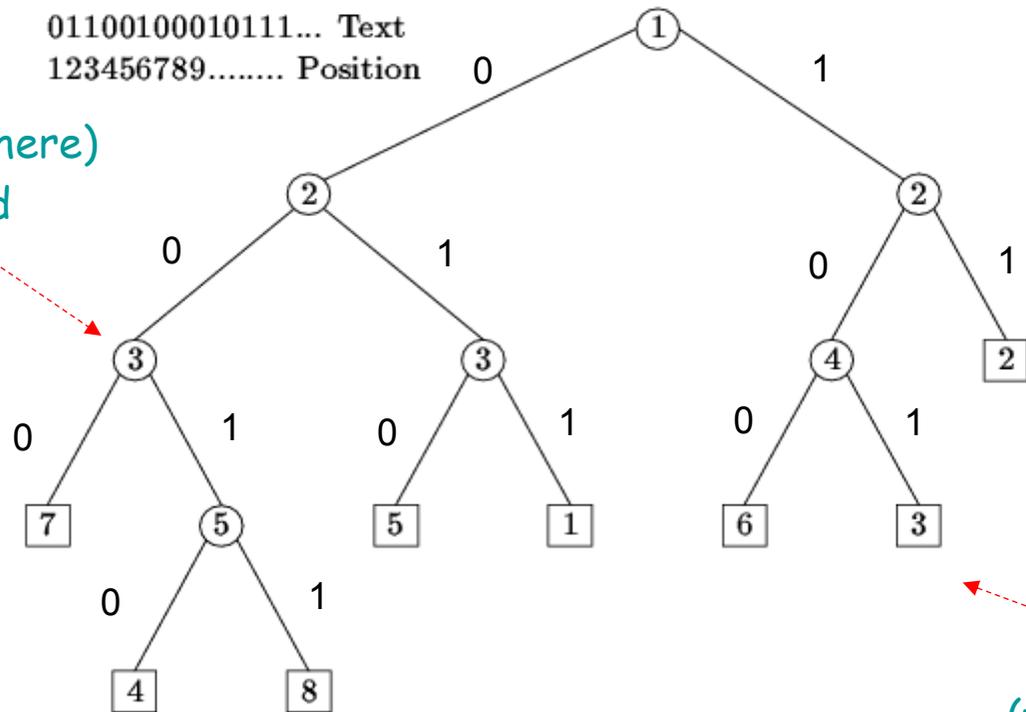
What if the query is "mo" or "modulation"?

# Suffix Trees (cont.)

- Another representation
  - PAT tree over a sequence of bits

The bit position of query used for comparison

- Absolute bit position (used here)
- Or the count of bits skipped (skip counter)



01100100010111... Text  
123456789..... Position

Pat tree when the sistrings 1 through 8 have been inserted.

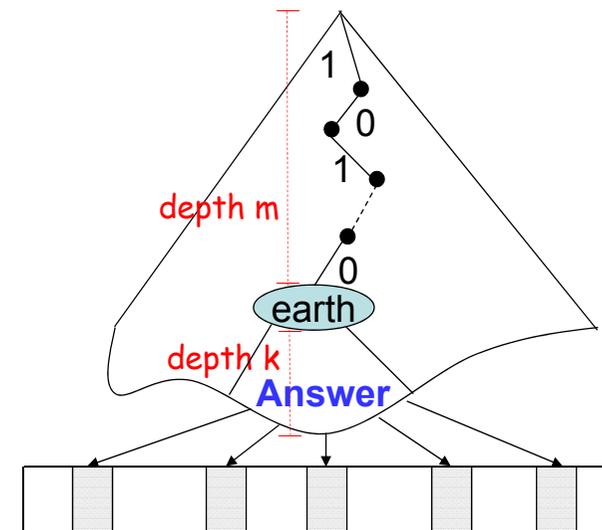
The key (text position)

Internal nodes with single descendants are eliminated!

Example query: 00101

# Suffix Trees: Search

- Prefix searching
  - Search the prefix in the tree up to the point where the prefix is exhausted or an external node reached  $O(m)$ ,  $m$  is the length in bits of the search pattern
  - Verification is needed
    - A single comparison of any of the sistrings in the subtree
  - If the comparison is successful, then all the sistrings in the subtree are the answer  $O(k \log k)$
  - The results may be further sorted by text order



# Suffix Trees: Search (cont.)

- Range searching
- Longest repetition searching
- Most significant or most frequent searching
  - Key-pattern/-word extraction

# Suffix Trees: Performance

- For a suffix tree built on 250 Mb of text
  - Single word or phrase (without supra-indices): 1 sec
  - Single word or phrase (with supra-indices): 0.3 sec

# Suffix Arrays

- **Basic Ideas**

- Provide the same functionality as suffix trees with much less space requirements
- The leaves of the suffix tree are traversed in left-to-right (or top-to-down here) order, i.e. lexicographical order, to put the points to the suffixes in the array
  - The space requirements is the same as inverted files
- Binary search performed on the array
  - Slow when array is large

$O(n)$ ,  $n$  is the size of indices

1	6	9	11	17	19	24	28	33	40	46	50	55	60
This is a text. A text has many words. Words are made from letters.													

Suffix array

60	50	28	19	11	40	33
----	----	----	----	----	----	----

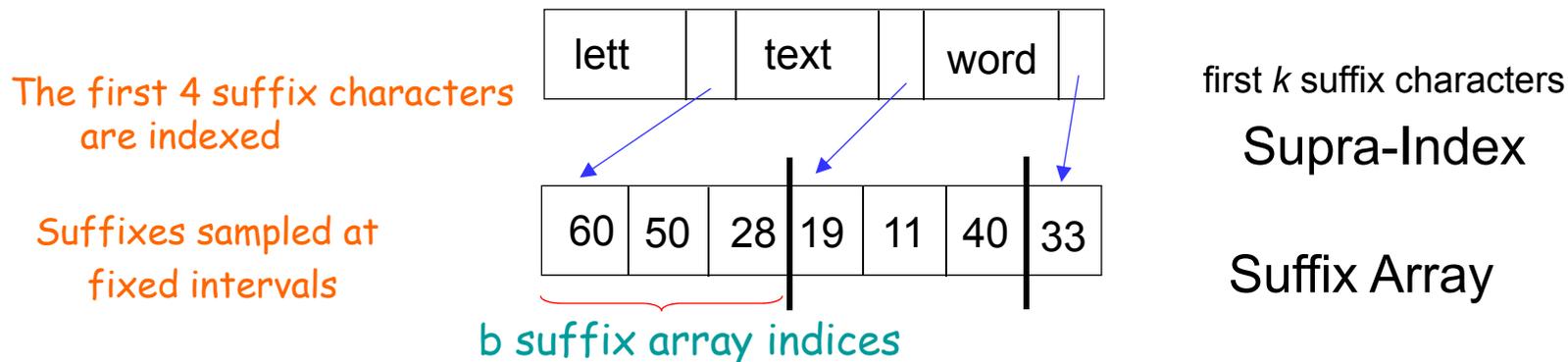
one pointer stored for each indexed suffix

(~40% overhead over the text size)

# Suffix Arrays: Supra indices

- Divide the array into blocks (may with variable length) and make a sampling of each block
  - Use the **first  $k$  suffix characters**
  - Use the **first word of suffix changes** (e.g., “text” (19) in the next example for nonuniformly sampling)
- Act as a first step of search to reduce external accesses (supra indices kept in memory!)

1	6	9	11	17	19	24	28	33	40	46	50	55	60
This is a text. A text has many words. Words are made from letters.													



# Suffix Arrays: Supra indices (cont.)

- Compare word (vocabulary) supra-index with inverted list

1	6	9	11	17	19	24	28	33	40	46	50	55	60
This is a <b>text</b> . A <b>text</b> has <b>many words</b> . <b>Words</b> are <b>made</b> from <b>letters</b> .													

letter	made	many	text	word
--------	------	------	------	------

Vocabulary  
Supra-Index

Suffixes sampled at  
fixed intervals

60	50	28	19	11	40	33
----	----	----	----	----	----	----

Suffix Array

60	50	28	11	19	33	40
----	----	----	----	----	----	----

Inverted List

- major difference {
- Word occurrences in suffix array are sorted lexicographically
  - Word occurrences in inverted list are sorted by text positions

# Suffix Trees and Suffix Arrays

- **Pros**

- Efficient to search more complex queries (phrases)
  - The query can be any substring of the text

- **Cons**

- Costly construction process
- Not suitable for approximate text search
- Results are not delivered in text position order, but in a lexicographical order